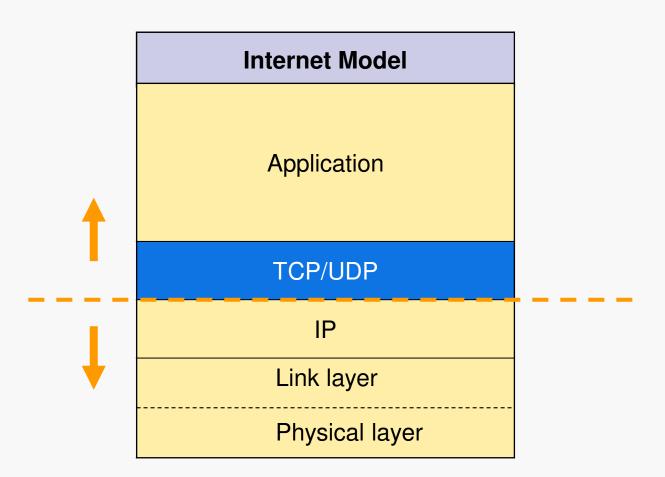
# **TCP/UDP Basics**



# **Transport Service Overview**

- Provide service to application layer by using the service provided by network layer
- Hide physical network
  - □ Hide processing complexity
  - □ Hide different network technologies and architectures
- Provide reliable, host-to-host transport

# Transport layer design issues

- Addressing
- Connection Establishment
- Connection Release
- Flow Control
- Error Detection and Crash Recovery

# Agenda

## TCP: Transmission control protocol (RFC 793)

- □ Addressing
- □ Connection Establishment
- □ Connection Release
- □ Flow Control
- □ Error Handling
- □ Interface and State Machine
- □ TCP application examples
- UDP: User datagram protocol (RFC 768)
  TCP vs. UDP

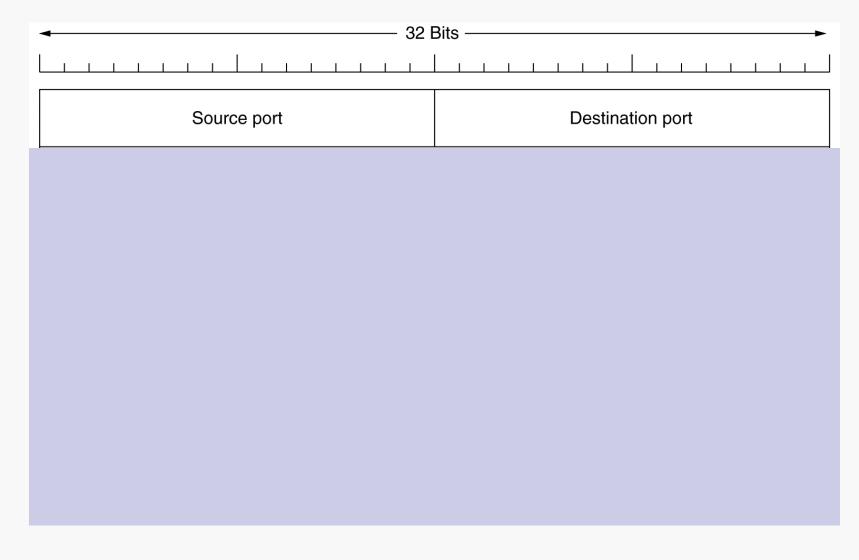
# **TCP** -- Addressing

- There are many network applications running on a host. When a packet arrive at network layer, how to know which application to send to?
  - $\square$  Port: there are  $2^{16} = 65536$  ports (0-65535) on one machine
  - □ One port is linked to only one application
  - □ One application may use many ports for different purposes
    - e.g. FTP: 20, 21

#### How a client knows which service uses which port?

- □ Permanent, well-known: often used service
  - 0-1023: well-known ports
  - 1024-49151: registered ports
  - 49152-65535: private ports
- Process server proxy and create service on-the-fly: temporary service
- □ Name server: for file service

# **TCP Addressing Header Fields**

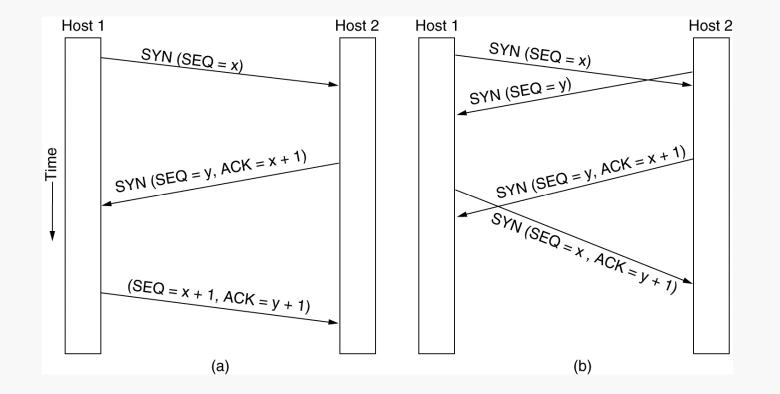


## TCP Connection Establishment – design issue

- Connection establishment becomes tricky when the network lose, delay and duplicate packets
   Bank example
- How to differentiate a new packet from a delayed, duplicated packet
  - □ Sequence number
    - Sequence number increase for each packet
    - Sequence number space issue:
      - □ Sequence number wrap back
        - A packet should avoid using a sequence number that another packet is using
  - □ A duplicated or delayed packet should die after a while
    - IP layer already handles this issue by 'Time To Live' header field

# **TCP Connection Related Header Fields**

# TCP Connection Establishment – solution three way handshake



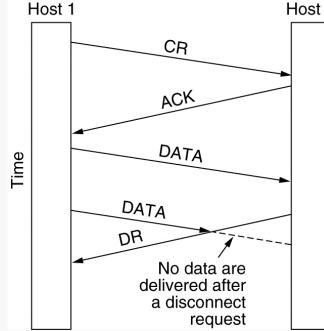
(a) TCP connection establishment in the normal case.

(b) Call collision.

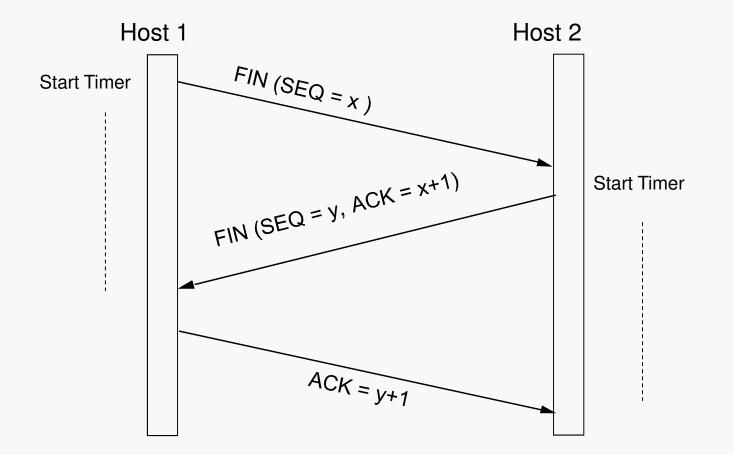
# TCP Connection Release – design options and issue (1)

Two release method: asymmetric and symmetric

Asymmetric release issue: possibility of losing data
Host 1



## TCP Connection Release – solution Three way handshake + timeout

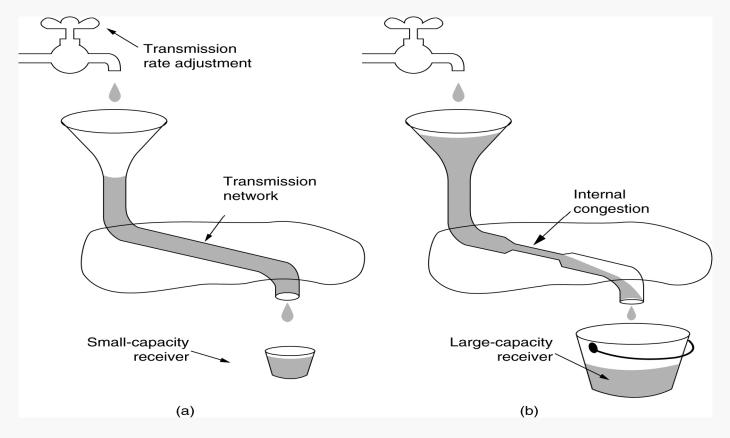


# TCP Flow Control – design issue

- Speed of data sending is critical
- □ Too fast:
  - network congestion or
  - receiving side overload

#### □ Too slow – type example

 waste of network resource or receiving memory

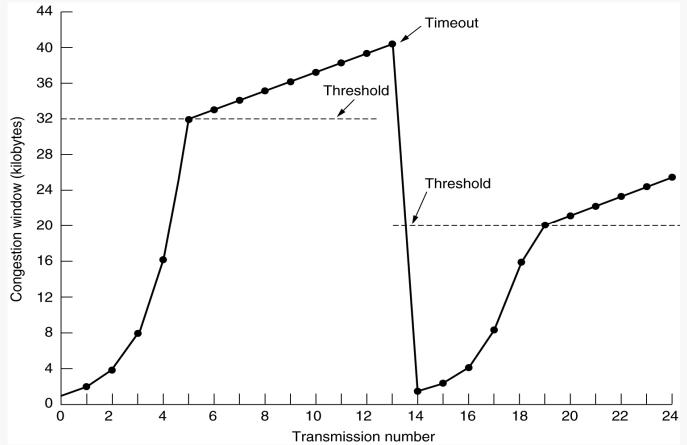


# TCP Flow Control – solution (1)

◄ 32 Bits →									
Source port							Destination port		
Sequence number									
Acknowledgement number									
TCP header length			A C K	P S H		Y	F I N	Window size	

- Windows maintained by both sending and receiving hosts
- Receiving side window size is decided by the available capability of receiving host's
- Sender maintains two windows
  - □ receiver window (got from receiving host), congestion window (to calculate)

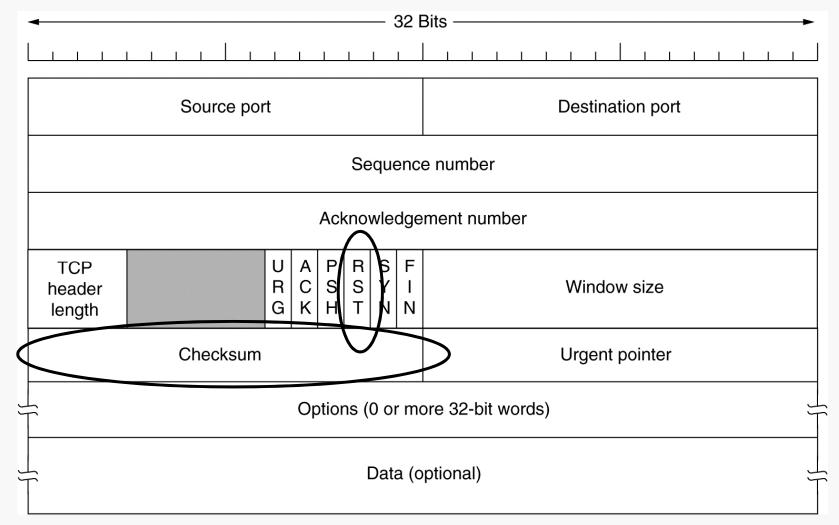
# TCP Flow Control – congestion window size calculation



slow start, reaching threshold, linear increment till timeout, recalculate threshold, slow start... till reaching receiver window size or timeout...

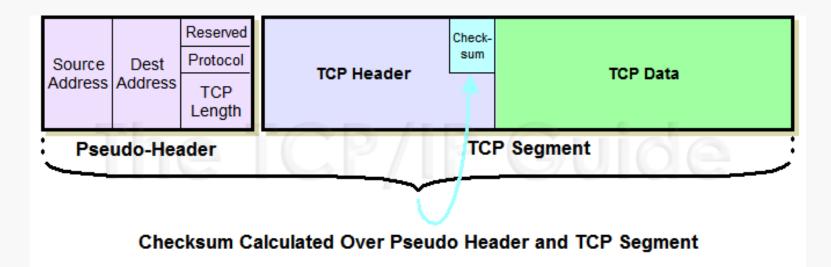
# **TCP Error Handling**

- Host crash and recovery
- Data error during transmission



# TCP Error Handling – TCP checksum

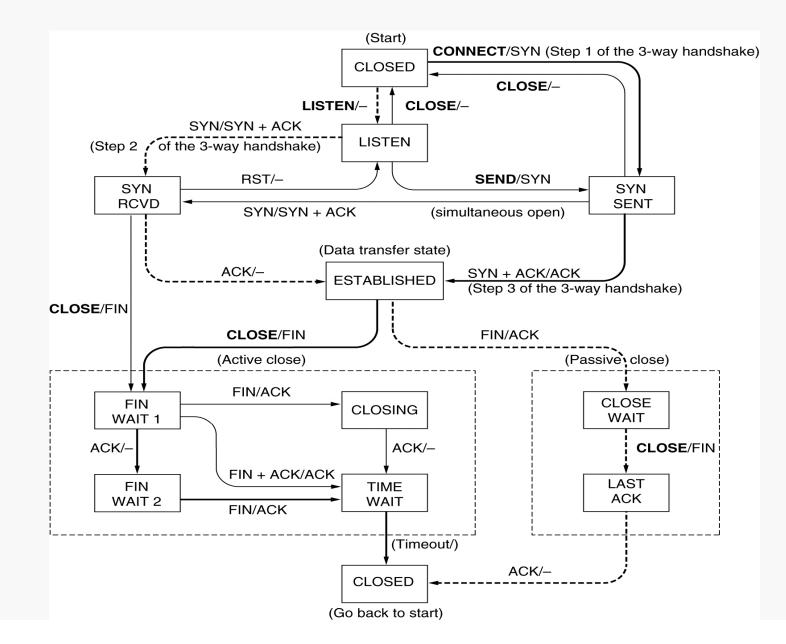
<b>←</b>	32 Bits	· · · · · · · · · · · · · · · · · · ·								
Source address										
Destination address										
0 0 0 0 0 0 0 0 Protocol = 6 TCP segment length										



# TCP service primitives

Primitive	Meaning						
SOCKET	Create a new communication end point						
BIND	Attach a local address to a socket						
LISTEN	Announce willingness to accept connections; give queue size						
ACCEPT	Block the caller until a connection attempt arrives						
CONNECT	Actively attempt to establish a connection						
SEND	Send some data over the connection						
RECEIVE	Receive some data from the connection						
CLOSE	Release the connection						

## **TCP** Finite State Machine



## **TCP** application examples

When to use TCP:

□ When an application need a reliable transport

### Examples

File Transfer Protocol : FTP (21)
Secure Shell: SSH (22)
Teletype Network: TELNET (23)
Simple Mail Transfer Protocol: SMTP (25)
Hypertext Transfer Protocol: HTTP (80)

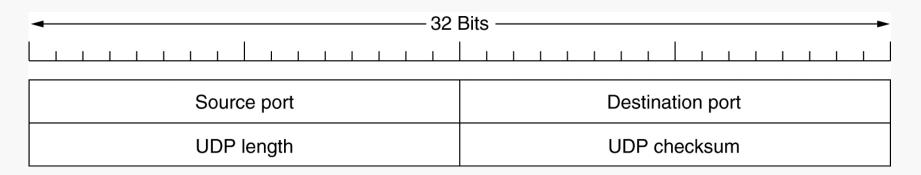
# Agenda

## TCP: Transmission control protocol (RFC 793)

## UDP: User datagram protocol (RFC 768)

- □ UDP header
- □ UDP properties
- □ UDP application examples
- TCP vs. UDP

# UDP Header



- UDP Destination Port: identifies destination process
- UDP Source Port: optional identifies source process for replies, or zero
- Message Length: length of datagram in bytes, including header and data
- Checksum: optional -- 16-bit checksum over header and data, or zero

# **UDP** Properties

- UDP provides an unreliable datagram service
   Packets may be lost or delivered out of order
   Message split into datagrams, user sends datagrams as packets on network layer
   No buffer at either sending or receiving side
   Unreliable but fast
   Full duplex
  - □ Application must deal with lost packets

# **UDP** Application Examples

### When to use UDP

- □ Reduce the requirement of computer resources
- The checking scheme has provided completely by the application program
- □ When using the Multicast or Broadcast to transfer
- □ The transmission of Real-time packets

#### Examples

- □ Trivial File Transfer Protocol , TFTP
- □ Simple Network Management Protocol , SNMP
- □ Dynamic Host Configuration Protocol , DHCP
- □ Domain Name System , DNS
- □ Routing Information Protocol , RIP
- □ Real-Time Transport Protocol , RTP

## TCP vs. UDP

TCPUDPconnection-orientedconnectionlessconfirmed serviceunconfirmed servicehigh overheadlow overhead(header 20 bytes)(header 8 bytes)flow controlno flow control

## References

## A. Tanenbaum, Computer Networks, Fourth Edition, Prentice Hall, 2003