An Anomaly Detection System based on Ensemble of Detectors with Effective Pruning Techniques

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ABSTRACT

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Anomaly detection systems are important tools for security. Unlike signature-based systems, anomaly detection can be used to detect new attacks for which signatures are now available.

To this end, anomaly detection techniques rely on machine learning techniques to model the normal behaviour of the system. This model is used as a baseline for the detection of anomalies during system operation.

The problem is that there is no one machine learning technique that can provide good accuracy. What we need is to combine multiple techniques. This is because ensemble methods have been used to improve the overall detection accuracy in traditional machine learning.

The combination consists of combining the outputs of several accurate and diverse models. To reduce the number of combination, and hence improve the efficiency of combination, in this thesis, we propose PBC (Pruning Boolean Combination), an efficient approach for selecting and combining anomaly detectors. PBC relies on two novel pruning techniques that we have developed to prune redundant and trivial detectors. Compared to existing work, PBC reduces significantly the number of detectors to combine, while keeping similar accuracy. We show the effectiveness of PBC when applying it to benchmarks data sets.

Much of the content of this thesis is adapted and expanded from a paper published at the 2015 IEEE International Conference on Software Quality, Reliability and Security (QRS). QRS is a merger of the SERE conference (IEEE International Conference on Software Security and Reliability) and the QSIC conference (IEEE International Conference on Quality Software) [1]. Dedicated to

My Family

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Chapter 1

Introduction

Intrusion Detection Systems (IDSs) have become important tools that help security administrators identify and analyze unauthorized computer or network activities, from both outsider and insider attacks. The dominant detection methodologies are signature-based and anomaly-based approaches. Signature-based (or misuse) IDSs look for events that correspond to the patterns of known attacks. They can provide a high level of accuracy, however, are limited to only detecting known attacks and are vulnerable to polymorphic ones (which are capable of changing their signatures as they propagate).

Anomaly detection systems (ADSs), on the other hand, monitor for significant deviations from normal system behavior. They are typically trained (using machine learning and data mining techniques) on data sets collected over a period of attack-free activities to learn normal system behavior, and then deployed to detect deviations from the expected system behavior. These deviations are reported as anomalous events, although they are not necessarily malicious activities or attacks because they may also correspond to coding or configuration errors. An ADS can detect novel attacks (not known during training time) but generate large number of false alarms, because it is difficult to obtain the complete descriptions of complex system behavior, which may change over time.

IDSs can also be categorized depending on their location of deployment into network- and host-based detection systems (a more comprehensive taxonomy of IDSs can be found in [3]). Network-based IDSs monitor and analyze traffic to and from all devices on the network. They may be located anywhere in the network; integrated in network devices (e.g., switches and routers) or as a standalone device. A host-based IDS runs on a host computer and monitors sensitive activities on this host, such as unauthorized access, modification of files, or system calls.

In this work, we focus on host-based anomaly detection using system calls – the gateway between user and kernel mode. Short sequences of system calls have been shown to be consistent with normal host operation, and can be used to detect attacks [4,5]. A large number of research studies have investigated different machine learning and data mining techniques for detecting anomalies in system call sequences [5–8].

Warrender and Forrest did one of the early work in this field [5]. They introduced the technique called Sequence Time-Delay Embedding (STIDE), which "acquires a model of normal behavior by segmenting training data into fixed-length sequence" [5]. This is done through sliding a detector window of length DW (hereafter DW) over the training data. Each of these sequences obtained from the data stream is stored in a "normal database" of sequences of length DW. Then a similarity metric is used to establish the degree of similarity between the test data and the model of normal behaviour that is obtained in the previous step. The sequence of length DW is obtained from the test data using a sliding window, and for each length DW sequence the similarity metric simply establishes whether that sequence exists or not in the normal database [9]. STIDE is a good starting point, but still has a number of issues. First, the best anomaly detection window size needs to be defined. Tan et al [9] showed that for the New Mexico data set the minimum window size should be 6, otherwise we are going to miss the detection of some intrusions. However, they mentioned in their work that a particular window size working for all data sets cannot be specified. Therefore, the minimum window size should be found based on the data in the training phase

The other issue with STIDE or other detectors that want to model the normal behaviour of the system, is that the learning phase is too time consuming, One should run the system in a very clean environment for a long time. Traces of the normal behaviour should be stored somewhere. And yet there is no guarantee that it is possible to model the normal behaviour of the system based on the traces that have been stored so far. In addition, the normal behaviour of the system will change in the course of time.

A study by Eskin et al [7] shows, that it is possible to estimate the proper window size for each data set instead of having a constant one. And by selecting and working on dynamic window sizes they get improved result. They estimate the optimal window size by two different techniques. The first method is Entropy modeling, a method which determines the optimal single window size, and the second, probability modeling which takes into account context dependent window sizes. In their work, they used a new method for modeling system call traces using sparse Markov transducers. By using sparse Markov transducers they were able to estimates the best window size depending on the specific system calls in the subsequence based on their performance over the training data [7].

One of the recent studies by Creech et al [7] is a good case that shows that using machine learning techniques on system call traces to find anomalies is still a hot topic. They proposed Extreme Learning Machine (ELM), which is a new and strong technique in the area of the neural network family. The advantage of ELM in comparison to previous techniques is its speed and scalability. Also, it is a very lightweight component, while still very accurate. However, based on the data there is an upper limit for its detection. In this thesis, we will address how we can use this technique alongside other techniques to increase the detection rate or decrease a false alarm.

Among all the techniques used to detect anomalies on system call traces, techniques based on Hidden Markov Models (HMMs) – probabilistic models for sequential data – have been shown to produce a high level of detection accuracy [5, 10–18].

The success of an ADS depends largely on the model of normal behavior. A single HMM may not, however, provide adequate approximation of the underlying data distribution of a complex host system behavior, due to the many local maxima of the likelihood function [16]. Ensemble methods have been used to improve the overall system accuracy by combining the outputs of several *accurate and diverse* models [19–22]. In particular, combining the outputs from multiple HMMs, each trained with a different number of states, in the Receiver Operating Characteristics (ROC) space according to the Iterative Boolean Combination (IBC), has been shown to provide a significant improvement in the detection accuracy of system call anomalies [23].

The IBC is a general decision-level combination technique that attempts to select the decision thresholds (from each input detector) and the Boolean functions that maximize the overall ROC convex hull of the combined ensemble [23]. As presented in Algorithm 1, given K soft detectors that assign scores or probabilities to the input samples, (which can be converted to a crisp detector by setting a threshold on the scores as further detailed in section 2.1), the IBC algorithm starts by combining the outputs of the first two detectors (in the order of input), and then proceeds *sequentially* by combining the resulting combinations with the outputs of third detector, and so on, until the K^{th} detector is combined. It can then re-iterate to combine the resulting combinations with the original detectors.

The IBC technique provides a practical way for combining relatively large number of detectors while avoiding the exponential explosion of Boolean combinations. As detailed in Section 3.3, applying all Boolean functions using an exhaustive bruteforce search to determine optimal combinations leads to an exponential number of combinations, which is prohibitive even for a small number of detectors [24]. Even the *pairwise* Bruteforce Boolean Combination (BBC2), which is used as a baseline reference in our experiments, requires an exponential number of combinations, which is equal to the square of the number of detectors (see Algorithm 2 without the pruning mechanisms).

However, the sequential combinations of soft detectors according to IBC still faces challenges, when a large number of detectors (K) is presented for combinations. First, it produces a sequence of combination rules that grows linearly with K(and the number of iterations), which becomes difficult to analyse and understand. Furthermore, it makes the algorithm sensitive to the order in which the detectors are input for combinations, which increases the effort required to find best subset for operations. In any case, combining all available detectors without any pruning mechanism may be inefficient due to the redundancy in their outputs.

In this work, we propose an ADS based on a Pruned Boolean Combination (PBC) algorithm, which employs two novel pruning techniques to select a subset of diverse and accurate detectors for combination, while discarding the remaining ones. Although both pruning techniques are based on Cohen's Kappa [25] measure, they differ in the way they compute the diversity and accuracy of the selected detectors for combinations. While the first technique relies only on Kappa measure (MinMax-Kappa), the second uses both Kappa and the ROC convex hull (ROCCH-Kappa). Therefore, the PBC approach proposed in this thesis provides an efficient way to prune and combine large number of detectors. This method avoids the exponential explosion of combinations of brute-force techniques, and reduces the sequence of combinations provided by IBC.

We evaluated the PBC-based ADS using ADFA Linux Data set (ADFA-LD)¹, which has been recently made publicly available on the website of the University of New South Wales [26]. The performance of the PBC using both pruning techniques are compared to that of IBC and BBC2 in terms of ROC analysis and time complexity. The results show that PBC with both pruning techniques are capable of maintaining similar overall accuracy as measured by the ROC curves to that of IBC and BBC2. However, the time required for searching and selecting (or pruning) the subset of detectors from the same ensemble of detectors is on average three magnitudes lower for PBC with MinMax-Kappa pruning than that of BBC2. Furthermore, the experimental results show that PBC with both pruning techniques always provides two crisp detectors for combination (during the operations), while IBC provided an average of 11 detectors to achieve the same operating point.

1.1 System Calls

A system call is a request in a Unix-like operating system made via a software interrupt by an active process for a service performed by the kernel. Users' requests are made from higher level applications to the kernel via a platform-dependent set of system calls. System calls provide an essential interface between a process and the operating system. Applications must invoke system calls to request kernel services, such as access to standard input/output, physical devices and network resources. Every system call has a unique number (known by the kernel) and a set of arguments. Some example of system calls used for file managements are as follows: open(), read(), write(), and close(). System call traces have different lengths that depending on the complexity and execution time of the process. Traces of a wide variety of attacks, including buffer overflows, symbolic link, decode and SYN floods may appear at the system call level and differ from normal behavior of privileged processes, ([4]; [27];

¹http://www.cybersecurity.unsw.adfa.edu.au/ADFA IDS Datasets/

[28]; [29]; [5]). Furthermore, after gaining root privileges on a host, attackers will typically try to maintain administrative access on the compromised host ([30]), by for instance installing backdoors and rootkits. Attackers will also attempt to distribute Trojan horses, using for instance comprised email addresses or shared network folders, to compromise other systems. These malicious activities may also invoke different sequences of system calls than those generated during normal execution of a process.

1.2 HMM-based Anomaly Detection using System Call Sequences

The temporal order of system calls is used to represent the normal behavior of a privileged process, while an unusual burst will occur during an attack [4,5]. In [4,5], the authors proposed a host-based ADS using a simple sequence matching technique for detecting anomalous system call sequences, generated from UNIX privileged processes [4]. System call traces provided for training are first segmented into fixed-length *contiguous* sequences, using a fixed-size sliding window, shifted by one symbol to build the normal profile. These sequences are then stored in a database that represents the normal process behavior. During operations, a sliding window (having the same size used to construct the normal profile) to scan the system calls generated by the monitored process for anomalies – sequences that are not found in the normal database.

Several statistical and machine learning techniques have been investigated over the last two decades for detecting system call anomalies [6]. Application of machine learning techniques include neural networks [31], k-nearest neighbors [32], Markov models [33, 34], Bayesian models [35]. Among these, techniques based on discrete HMMs have been shown to produce a high level of detection accuracy [5, 10–16, 36– 38]. A discrete HMM is a stochastic process for sequential data [39,40]. An HMM is determined by two interrelated mechanisms – a latent Markov chain having a finite number of states N, and a set of observation probability distributions, each one associated with a state. Starting from an initial state $S_i \in \{S_1, ..., S_N\}$, determined by the initial state probability distribution π_i , at each discrete-time instant, the process transits from state S_i to state S_j according to the transition probability distribution a_{ij} . The process then emits a symbol v_k , from a finite alphabet $V = \{v_1, ..., v_M\}$ of size M symbols, according to the discrete-output probability distribution $b_j(v_k)$ of the current state S_j . HMM is commonly parametrized by $\lambda = (\pi, A, B)$, where the vector $\pi = {\pi_i}$ is the initial state probability distribution, matrix $A = {a_{ij}}$ is the state transition probability distribution, and matrix $B = {b_j(v_k)}$ is the state output probability distribution, $(1 \le i, j \le N \text{ and } 1 \le k \le M)$.

A well trained HMM provides a compact detector that captures the underlying structure of a process based on the temporal order of system calls, and detects deviations from normal system call sequences with high accuracy and tolerance to noise. Training an HMM from a sequence (or a block) of observation symbols, $o_{1:T}$, aims at estimating HMM parameters λ to best fit the training data. Typically, parameters estimation consists of maximizing the likelihood of the training data over HMM parameters space, $P(o_{1:T} \mid \lambda)$. Since this likelihood depends on the latent states, there is no known analytical solution to the learning problem. Iterative optimization techniques, such as the Baum-Welch (BW) algorithm [41], are applied to estimate the HMM parameters over several training iterations, until the likelihood function is maximized. During operation, the likelihood of a new observation sequence $o_{1:T}$ given a trained HMM λ , $P(o_{1:T} \mid \lambda)$ is typically evaluated by using the Forward-Backward (FB) algorithm [39,40]. Setting a threshold on the output probabilities of HMM, provides a decision whether the sequence is normal or anomalous. Some researchers investigated the effect of the number of states on the performance of HMM detectors, and found that N value may have a great impact on the overall performance [16,42]. In particular, combining the outputs from multiple HMMs, each trained with a different N value, in the ROC space according to the IBC technique (as described in the next section), has been shown to provide a significant improvement in the detection accuracy of system call anomalies [23].

Chapter 2

Detector Combination

2.1 Boolean Combination of Detectors in the ROC Space

This section provides information about ROC analysis and summarizes the Boolean combination in the ROC space [23]. A crisp detector outputs a decision or a class label (e.g., normal or anomaly) while a soft detector such as HMM assigns scores to the input samples, which can be converted to a crisp detector by setting a decision threshold on the scores. Given the responses of a crisp detector on a validation set, the true positive rate (tpr) is the proportion of positives correctly classified over the total number of positive samples. The false positive rate (fpr) is the proportion of negative samples. The positive (or target) class is typically the class of interest, which is the anomalous class for an ADS.

A ROC curve is a plot of tpr against fpr [43]. A crisp detector produces a single data point in the ROC space, while a soft detector produces a ROC curve by varying the decision thresholds. In practice, an empirical ROC plot is obtained by connecting the observed (tpr, fpr) pairs of a soft detector at each decision threshold. A point a is

superior to another point b in the ROC space, if $fpr(a) \leq fpr(b)$ and $tpr(a) \geq tpr(b)$. The ROC convex hull (ROCCH) is therefore the outer envelope connecting superior points in the ROC space. A ROC curve allows to visualize the performance of detectors and to select optimal operational points, without committing to a single decision threshold or to fixed error costs. The area under the ROC curve (AUC), or under the ROCCH, provides a general measure for evaluation and selection of detectors [43].

The IBC is a general decision-level combination technique that attempts to select the decision thresholds (from each input detector) and the Boolean functions that maximize the overall ROCCH of the combined ensemble [23]. The core of IBC (only for the first iteration) is described in Algorithm 1.

The IBC applies each Boolean function to combine the responses corresponding to each decision threshold from the first detector to those from the second detector. Fused responses are then mapped to vertices in the ROC space, and their ROC convex hull (ROCCH) is computed. Vertices that are superior to the ROCCH of original detectors are then selected. The set (\$) of decision thresholds from each detector and Boolean functions corresponding to these vertices is stored, and the ROCCH is updated to include emerging vertices. The responses corresponding to each decision threshold from the third detector are then combined with the responses of each emerging vertex, and so on, until the last detector in the pool is combined. The BC technique yields a final ROCCH for visualization and selection of operating points, and the set of selected thresholds and Boolean functions, \$, for each vertex on the composite ROCCH to be applied during operations. Although not shown in Algorithm 1, the original IBC algorithm can iterate by re-combining the resulting combination (of all detectors) on the ROCCH with each of the original detectors (sequentially) until the ROCCH stops improving [23]. However, as stated previously, combining all available detectors without pruning may be inefficient due to the redundancy in their outputs. The sequential combinations of soft detectors, illustrated in the loop in line 16 of Algorithm 1, produces a sequence of combination rules that grows linearly with K (and the number of iterations), which becomes difficult to track, analyse and understand when the value of K becomes large. In addition, the IBC algorithm is sensitive to the order in which the detectors are input for combinations, which increases the effort required to find best subset for operations.

2.2 Synthetic Data Set

To better explain the reason for choosing the proper combination of detectors among millions of trained detectors, in this section, examples of training, pruning and combining detectors on a synthetic data set are discussed.

For this purpose, three known synthetic data sets are chosen [44].

- 1. Lithuanian
- 2. Banana (Moon)
- 3. Circle

Unlike the n-dimensional real data sets, it is possible to plot synthetic data sets and to illustrate the trained models generated by them. This would assist an individual to observe what happens when unnecessary trained models are pruned and only those that can increase the detection accuracy remain.

Figure 2.1, 2.2 and 2.3 shows the Lithuanian, Banana and Circle data sets. Figure 2.4, 2.5 and 2.6 show the related trained models on these data sets.

Another purpose of this example is to show that we can reach the same level of detection by combining very simple and naïve model detectors, instead of spending time to find the proper model and to train that model on the data set, which is



Figure 2.1: Lithuanian Data Set



Figure 2.2: Banana (Moon) Data Set



Figure 2.3: Circle Data Set

usually a time consuming task. In this example, hundreds of linear detectors are trained the data set. In order to have diverse learner detectors, we changed the ratio of data and their label during training to obtain as many diverse detectors as possible.

Another purpose of this example is to show that we can reach the same level of detection by combining simple and naïve model detectors, instead of spending time to find the proper model and to train that model on the data set. In this example, hundreds of linear detectors are trained on data set. In order to have diverse learner detectors, the ratio of data and their label during the training process are changed to obtain as many diverse detectors as possible.

Figure 2.4, 2.5 and 2.6 show different trained detectors on these synthetic data sets respectively. None of them, however, can detect all the classes. Yet by combining these simple detectors, a better data set classification is achieved as it is shown later.



Figure 2.4: Linear detectors trained on the Lithuanian Data Set



Figure 2.5: Linear detectors trained on the Banana (Moon) Data Set



Figure 2.6: Linear detectors trained on the Circle Data Set

By selecting the right models and combining them the detection rate ma increase, as shown in Figures 2.8, 2.7 and 2.9. The figures on the left show the original detectors and their ROC before combination. After selecting the proper detectors and combining them together (i.e. after just one iteration) the accuracy portrayed on the right figure is achievable, which show the combined detectors (in yellow) and their improved ROC curve. The next step is to show whether combination provides a better result or not, and how to obtain the same result just by combining parts of the trained models.



Figure 2.7: ROC comparison on Lithuanian data set before and after first round of combination



(a) ROC before combination

(b) ROC after combination

Figure 2.8: ROC comparison on Banana data set before and after first round of combination



Figure 2.9: ROC comparison on Circle data set before and after first round of combination

For each data set, it is concluded there is no single detector that can classify the data set properly.

As it is shown in Figure 2.10, with just one Linear Discriminant Analysis (LDA), it is not possible to classify the entire classes with 100% accuracy. To obtain a more accurate result, however, two of these detectors may be chosen and combined. As shown in Figure 2.11, a detector from Figure 2.10a and detectors from 2.10d are selected and combined.



Figure 2.10: Some of best LDAs (Linear Discriminant Analysis) that can be trained on the Lithuanian data set



Figure 2.11: Selected LDA and their respective point compare to current ROC

If these two detectors with an *AND* operation are combined, as shown in Figure 2.12, , a more reasonable result is obtained. The red point is the mapping of the combined detectors. If new ROC based on this new generated point will be recalculated, a greater AUC is achieved.



Figure 2.12: Selected detectors and the respective combined point on the ROC curve

A similar approach can be applied to the Banana data set. As it is shown in Figure 2.13, like the Lithuania data set, it is not possible to classify the data set completely with one detector. In this specific example, at least three detectors are needed to precisely classify the data in an accurate way.

Given the fact that three detectors are needed for data set classification, combining the selected detectors shown in Figure 2.13a, Figure 2.13c and Figure 2.13f may result in achieving the highest accuracy among all other detectors. Figure 2.14 portraits the whole mentioned detectors in one plot



Figure 2.13: Some of best LDA (Linear Discriminant Analysis) that can be trained on Banana data set



Figure 2.14: Best potential selected detectors on Banana data set

These examples show that it is not always necessary to train complex models, which takes time and resources, over the data set to obtain the best result possible. The same result can be achieved by combining two or more simple models, which can be trained much faster. Also by increasing the number these simple models, most of the cases (including the outliers) in the data set may be covered.

Moreover, due to the imperfection nature of likely every model, training the complex models would not guarantee the best achievable result. Therefore, a combination is needed to enhance the outcome. Since it is not possible to train and to combine thousands of models, our approach on how to select a set of meaningful models to improve the expected result will be shown in the next chapter.

The Circle data set is a good example where combining in an iterative way improves the accuracy shown in Figure 2.15. W. Since it is not possible to combine any two detectors which could classify the whole data set, the entire detectors are selected and combined two by two iteratively. The result of classifying iteratively


the whole data set with the help of a few simple detectors is shown in Figure $\ 2.16.$

Figure 2.15: Series of 2 selected LDA (Linear Discriminant Analysis) that can be trained and combined on the Circle data set



Figure 2.16: Best potential selected detectors on the Circle data set

Al	gorithm 1: $IBC(D_1, D_2,, D_K, \mathcal{V})$: Iterative Boolean Combination				
	input : K soft detectors $(D_1, D_2,, D_K)$ and a validation set \mathcal{V} of size $ \mathcal{V} $				
	output: ROCCH of combined detectors.				
	- Each vertex is the result of 2 to K combination of crisp detectors.				
	- Each combination selects the best decision thresholds from different				
	detectors (D_i, t_j) and Boolean function (stored in the set \mathcal{S})				
1	$n_k \leftarrow \text{number of decision thresholds of } D_k \text{ using } \mathcal{V} // \text{ num. of vertices on ROC(D}_k).$				
2	$BooleanFunctions \leftarrow \{a \land b, \neg a \land b, a \land \neg b, \neg (a \land b), a \lor b, \neg a \lor b, a \lor \neg b, \neg (a \lor b), a \oplus b, a \equiv b\}$				
3	compute ROCCH ₁ of the first two detectors $(D_1 \text{ and } D_2)$				
4	allocate F an array of size: $[2, n_1 \times n_2]$ // temporary storage of combination results.				
5	for each $bf \in BooleanFunctions$ do				
6	$\mathbf{for}\ i \leftarrow 1 \ \mathbf{to}\ n_1 \ \mathbf{do}$				
7	$R_1 \leftarrow (D_1, t_i)$ // responses of D_1 at decision threshold t_i using V .				
8	for $j \leftarrow 1$ to n_2 do				
9	$R_2 \leftarrow (D_2, t_j) \qquad \qquad // \text{ responses of } \mathbf{D}_2 \text{ at decision threshold } \mathbf{t_j} \text{ using } \mathcal{V}.$				
10	$R_c \leftarrow bf(R_1,R_2)$ // combine responses using current Boolean func.				
11	$\textbf{compute } (tpr, fpr) \text{ of } \boldsymbol{R_c} \text{ using } \mathcal{V} \qquad // \text{ map combination to ROC plane}$				
12	push (tpr, fpr) onto F				
19	$ \sqsubseteq $				
14	$n_{rrr} \leftarrow \text{number of emerging vertices}$				
15	$S_2 \leftarrow \{(D_1, t_i), (D_2, t_i), bf\}$ // set of selected decision thresholds from each detector and Boolean functions for emerging				
10					
10	$ \begin{array}{c} 16 10 7 6 5 10 7 10 $				
10	anotate F of size. $[2, n_k \times n_{ev}]$				
10	foreach $bf \in BooleanFunctions$ do				
19	$\begin{bmatrix} \mathbf{R}_{i} \\ \mathbf{R}_{i$				
20	for $i \in J_{k-1}(i)$ // responses from previous combinations.				
41 22	$\begin{vmatrix} \mathbf{R} \\ \mathbf{R} $				
22 23	$ \begin{array}{c} R_{k} \leftarrow (D_{k}, v_{j}) \\ R_{k} \leftarrow bf(R, R_{k}) \end{array} $				
20	compute (trr, frr) of \mathbf{R}_{\star} using \mathcal{V}				
25	$\mathbf{push} (trr frr) \text{ onto } \mathbf{F}$				
_0					
26	compute $ROCCH_k$ of all ROC points in F				
27	$n_{ev} \leftarrow$ number of emerging vertices				
28	$\mathbb{S}_k \leftarrow \{\mathbb{S}_{k-1}(i), (D_k, t_j), bf\}$				
	$//S_k$ is the set of the selected subsets from the previous combinations; the decision thresholds from the newly-combined detector;				
	and the Boolean functions that yields to the emerging vertices on the ROCCH.				
29	- 29 store $S_k : 2 < k < K$				
30) return $ROCCH_K$				
	**				

Algorithm 2: $PBC(D_1, D_2, ..., D_K, \mathcal{V})$: Pruned Boolean Combination **input** : K soft detectors (D_1, D_2, \dots, D_K) and a validation set \mathcal{V} of size $|\mathcal{V}|$ output: ROCCH of combined detectors. - Each vertex is the result of exact 2 combination of crisp detectors. - Each combination selects the best decision thresholds from different detectors (D_i, t_j) and Boolean function (stored in the set S) $\mathbf{1} \quad n_k \leftarrow \text{number of decision thresholds of } D_k \text{ using } \mathcal{V}$ // num. of vertices on $ROC(D_k)$. 2 let $n = \sum_{k=1}^{K} n_k$ **3** BooleanFunctions $\leftarrow \{a \land b, \neg a \land b, a \land \neg b, \neg(a \land b), a \lor b, \neg a \lor b, a \lor \neg b, \neg(a \lor b), a \oplus b, a \equiv b\}$ 4 allocate C an array of size: $[|\mathcal{V}|, n]$ // storage of all crisp detectors' decisons. 5 convert soft detectors to crisp detectors 6 for $i \leftarrow 1$ to K do for $j \leftarrow 1$ to n_i do 7 $\boldsymbol{R} \leftarrow (D_i, t_j)$ 8 // responses of \mathbf{D}_i at decision threshold \mathbf{t}_j using $\mathcal{V}.$ push R onto C9 10 **choose** Pruning Technique {MinMax-Kappa, ROCCH-Kappa} **reduce** n to U// $U \ll n$: is a user defined max number of detectors 11 12 return $\mathcal{C}_{ ext{selected}} \leftarrow C$ - Pruned Detectors // Subset of size U detectors selected from all original detectors and returned for combination 13 allocate F an array of size: $[2, U^2 \times size(BooleanFunctions)]$ // temporary storage of combination results. 14 foreach $bf \in BooleanFunctions$ do 15 for $i \leftarrow 1$ to U do $R_1 \leftarrow \mathfrak{C}_{ ext{selected}}[i]$ 16 // Retrieve Decision Vector for $j \leftarrow 1$ to U do $\mathbf{17}$ $R_2 \leftarrow \mathfrak{C}_{ ext{selected}}[j]$ $\mathbf{18}$ $R_{c} \leftarrow bf(R_{1}, R_{2})$ 19 // combine responses using current Boolean func. **compute** (tpr, fpr) of $\mathbf{R_c}$ using \mathcal{V} // map combination to ROC plane 20 21 push (tpr, fpr) onto F22 compute ROCCH of all ROC points in F **23** $n_{ev} \leftarrow$ number of emerging vertices **24** $\& \in \{(D_1, t_i), (D_2, t_j), \dots, (D_k, t_k), bf\}$ // set of selected decision thresholds from each detector and Boolean functions for emerging vertices. 25 store S; return ROCCH

Chapter 3

Pruning

3.1 Introduction

Pruning is a very popular practice for decreasing the cost and time of combining multiple detectors. There are two main challenges when it comes to pruning. The first one is to determine where pruning should be applied (e.g., feature level, model level, etc.) and the second challenge is to determine the pruning criteria.

Generally speaking, pruning can be done at four different levels: sensor level, feature level, matching score level, and decision level [2], as illustrated in Figure 3.1

In this thesis, the main focus of pruning detectors is at the decision level. However, there are some scenarios that requires pruning at the matching score level. For example, when there are thousands of detectors; it is not beneficial to convert them into a few response vectors and combine everything together. The reason is that we will have too many redundant response vectors, which hinders the efficiency of the pruning process.

After seeing different levels where pruning can be applied, the question is how to prune the pool of available detectors and make a basket of reasonable detectors. To have a reasonable basket of remaining detectors, two criteria should be considered.



Figure 3.1: Different levels of combination or pruning: sensor level, feature level, matching score level, and decision level [2].

Firstly, we have to define some accurate detectors in the basket, with this we guarantee that the minimum accuracy we obtain after pruning and the combination is not worse than any of the detectors we had in the pool. Also, by choosing and combining the accurate ones, we increase the probability of getting better results than combining two random or naïve detectors.

The second criteria is the diversity in the basket of detectors. If very similar detectors are selected in the basket, we should not expect to see a major improvement after combining. One way to increase the diversity of the basket is to try to find and select detectors that complement the errors of already selected detectors. Since we have already selected accurate detectors in the basket, it is a good practice to choose the detectors that complement the errors these detectors exhibit.

To see the importance of diversity and how it can help to improve the result of the combination, we show in the following example, the importance of selecting the initial detectors and respective detectors compare to them especially, the ones that tries to complement first detector's errors. By combining these two detectors with a right Boolean operator we will reach a very promising result. As it is shown in Figure 3.2, among all the trained detectors, we only need to find a way to select one detector that is accurate, and in Figure 3.3, we demonstrate how to select a detector, that is very diverse to previously selected ones in order to increase the chance of getting better results after combining. The two selected detectors in Figure 3.2 and Figure 3.3 are completing each other errors.



Figure 3.2: Selecting accurate detectors



Figure 3.3: Selecting diverse detector compares to previously selected

3.2 Pruning Boolean Combination (PBC) Approach

In this section, we describe our Pruned Boolean Combination (PBC) algorithm, which is based on two novel pruning techniques that select a subset of diverse and accurate detectors for combination, while discarding the remaining ones, this happened at decision level. PBC can avoid both the exponential explosion of combinations provided by the brute-force approach and the sequential combination of IBC. The main difference is that the PBC algorithm proceeds by thresholding all available soft detectors into crisp ones. This step is, in fact, not essential for PBC because thresholding could be done outside the algorithm and the crisp detectors are directly input for pruning. In contrast with IBC, all available crisp detectors are input to one of the proposed pruning techniques, MinMax-Kappa or ROCCH-Kappa described in Section 3.2.3 and 3.2.4, to select the best subset of crisp detectors for Boolean combination and control its size U (i.e., the number of combined detectors). Without pruning brute-force pairwise combination of all available detectors may not be feasible for large number of detectors, $O(N^2)$ (where N is number of available detectors), for further details see Section 3.3.

After seeing the main pruning idea in Section 3.2.3 and 3.2.4 we introduce a different distance correlation (measure of agreement between vectors or dependence coefficient) 3.2.2 rather than kappa to show that the results of PBC are consistent, and as long as a proper measure is used we can get promising result either using it with MinMax- or ROCCH-. In Section 3.2.5 we explained how we used the MinMax and ROCCH technique again and just changed the measure of agreement from kappa to dCor or dCov to prune the pool of detectors.

The boost in performance achieved with an ensemble of detectors is often attributed to the concept of diversity. While it is generally accepted that an ensemble should contain diverse models to improve the performance, there is no clear definition of diversity neither a consensus about the measure of diversity and its computation. In practice, several measures of diversity have been proposed to quantify the level of agreement or measure the dependency among the ensemble members [45].

3.2.1 Kappa Measure

Cohen's Kappa statistic (or simply Kappa hereafter) is one of the well-known and widely used measure of agreement between raters [25]. Kappa has lately gained some popularity for ensemble combination, especially the Kappa-error diagrams which help visualizing individual accuracy and diversity in a two dimensional plot [46, 47]. Our pruning techniques, described in the next sections, are based on Kappa and inspired by the Kappa-error diagram only for visualization of Kappa against the false positives and true positives (as shown in Figures 3.4 and 3.5).

	D_2 correct	D_2 wrong
D_1 correct	a	b
D_1 wrong	с	d

Table 3.1: Contingency table between two detectors' decisions

Consider the contingency table of two detectors, D_1 and D_2 , presented in Table 3.1, where, for instance, a is the number of examples on which both detectors agree. The sum of all element in Table 3.1 is equal to the size of validation set, $a + b + c + d = |\mathcal{V}|$. The Kappa (κ) measure between two detectors is therefore computed based on the element of the contingency table according to Equation 3.1.

$$\kappa = \frac{2(ad - bc)}{(a+b)(b+d) + (a+c)(c+d)}$$
(3.1)

Kappa takes on values between -1 and 1; lower values means a high level of disagreement or more diverse opinions, while higher values indicate a high level of agreement or similarity in responses between detectors. When both detectors provide the same vector of decisions (they agree on every example) then $\kappa = 1$. On the other hand, when $\kappa = 0$ the detectors are independent (any agreement is totally due to chance). Negative Kappa values can be interpreted as both detectors agrees less than what would be expected just by chance. More importantly, negative values account for negative correlations, which can be useful for combination, but this rarely occurs in practice [46].

3.2.2 Distance Covariance (dCov) and Distance Correlation (dCor) Measure

Distance correlation is a measure of statistical dependence between two vectors of arbitrary. The measure of dependence is zero if and only if the random variables are statistically independent. Distance correlation is a new measure of dependence between random vectors introduced by [48].

In order to understand the definition of distance covariance and distance correlation we should first start with distance covariance.

Assume we have two random vectors, the vectors are X_n and Y_n , k = 1, 2, ..., n. Firstly we should compute all pairwise distances:

$$a_{j,k} = ||X_j - X_k||$$
$$b_{j,k} = ||Y_j - Y_k||$$

By this, we computed the n by n distance matrices (aj, k) and (bj, k). After this we need to calculate all doubly centered distances:

$$A_{j,k} := a_{j,k} - \bar{a}_{j.} - \bar{a}_{.k} + \bar{a}_{..} \tag{3.2}$$

$$B_{j,k} := b_{j,k} - \bar{b}_{j,} - \bar{b}_{.k} + \bar{b}_{..}$$
(3.3)

In equation 3.2 $\bar{a}_{j.}$ is the j-th row mean, $\bar{a}_{.k}$ is the k-th column mean, and $\bar{a}_{..}$ is the grand mean of the distance matrix of the X sample. The exact same notation applies on b too. One thing to note is that in the matrices of centered distances $A_{j,k}$ and $B_{j,k}$ sum of all rows and all columns equal to zero.

By having all of these value we can calculate the distance covariance like this:

$$dCov_n^2(X,Y) := \frac{1}{n^2} \sum_{j,k=1}^n A_{j,k} B_{j,k}$$
(3.4)

Interesting point about distance covariance is that we can define it with help of Pearson's covariance.

$$dCov^{2}(X,Y) := cov(||X - X'||, ||Y - Y'||) - 2cov(||X - X'||, ||Y - Y''||)$$
(3.5)

After calculating the distance covariance we can use it to calculate distance correlation, The distance correlation of two random vectors is calculated by dividing their distance covariance by the product of their distance standard deviations.

$$dCor(X,Y) := \frac{dCov(X,Y)}{\sqrt{dVar(X)dVar(Y)}}$$
(3.6)

3.2.3 MinMax-Kappa Pruning

The proposed pruning technique starts by computing the Kappa values between each detector's decision vector and the true decision labels (or ground truth), and then sorting them in ascending order. Detectors with large Kappa values ($\kappa \approx \kappa_{max}$) are accurate and hence should be selected; however, they provide less diverse decisions among themselves. Therefore, the technique attempts to select complementary detectors by choosing those with Kappa values close to $\kappa \approx \kappa_{min}$. In practice, κ_{max} and κ_{min} values depend on the data and the detectors. Trivial detectors (providing always either positive or negative decisions) also reside at $\kappa_{min} \approx 0$. In such cases, these are filtered out before selected the complementary detectors. When detectors with negative Kappa values exist, it is always beneficial to select them since they provide detectors with negative correlations or complementary errors compared those with κ_{max} .

The MinMax-Kappa pruning technique takes two parameters, the total number of detectors and the ratio of detectors to be selected close to κ_{max} and κ_{min} . In our experiments, we experimented with different parameters sets and presented the results for 50 selected detectors with a ration of 50% (which means half of the detectors are selected from the region with higher values of Kappa, while the remain half are chosen from the region with lower values). In fact, these are user-defined parameters that are constrained by the resources available during operations. Our experiments showed that the results of the pruning algorithms are not very sensitive to small changes in these parameters

Figure 3.4 shows an example of the selected detectors from our experiment, according to the MinMax-Kappa technique. The Kappa values (on the X-axis) for the same selected detectors are plotted against false positive (Figure 3.4a) and true positive (Figure 3.4b) rates. These points (large, blue) are the selected detectors, $C_{selected}$ in Algorithm 2. All remaining detectors (small points) are pruned. For illustration, Figure 3.6a map the selected point to the ROC space, which could also be compared to our second pruning technique.

3.2.4 ROCCH-Kappa Pruning

This technique also tries to select accurate detectors and then find a set of detectors with complimentary errors. In contrast with MinMax-Kappa, this technique considers the detectors that are on the facet of the ROCCH in the ROC space. These detectors are selected since they are the most accurate available detectors that covers the whole range of *fpr* and *tpr* trade-off. The Kappa measure is then used to select diverse detectors with reference to those on ROCCH. The technique computes the Kappa value of each detector on the ROCCH and all the remaining detectors, sort the resulting values in ascending order. Detectors with larger values of Kappa are discarded, since they provide similar decisions (or errors) to those provided by the select detectors (on the ROCCH).

The most diverse detectors are those that provide negative or close to zero Kappa values ($\kappa \leq 0$), however are not trivial detectors (providing always either positive or negative decisions). The only parameter of the ROCCH-Kappa is the number of detectors to be selected for each detector on the ROCCH. In practice, the number of detectors on the ROCCH is small ([3-20] detectors). Figure 3.5 shows an example of the selected detectors from our experiment, according to the ROCCH-Kappa technique, where the Kappa values (on the X-axis) for the selected detectors are plotted against false positive (Figure 3.5a) and true positive (Figure 3.5b) rates. Similarly, these selected points (large, blue) are the $C_{selected}$ in Algorithm 2, while all remaining detectors (small points) can be pruned. Figure 3.6b map the detectors selected by ROCCH-Kappa to the ROC space, which could be compared to those selected by MinMax-Kappa in Figure 3.6a.

3.2.5 MinMax-, ROCCH- Pruning with dCor and dCov

These techniques are using the logic we proposed in section 3.2.3 and section 3.2.4. The novelty behind them is to use newer and stronger distance correlation, which is described in section 3.2.2. With the help of these two new metrics and integrating them with MinMax and ROCCH techniques; we are more confident in the performance of our PBC method. With this approach the newly proposed PBC method is working well and is stable at all times. More specifically because we use this method with a rational measure of agreement, we can expect to get the reasonable results.

Here in both dCor and dCov; if the result of comparing two different vectors is equal to 0 indicates that the vectors are surely independent of each other; and if the result is equal to 1, it means they are exactly similar to each other.

3.2.6 Expected Value of Boolean Combination

Expected Value of Boolean Combination (EVBC) Our other approach for pruning classifiers is using Expected value. To do that we calculate expected values of combinations of each classier with every other classifiers. This is an iterative method that would be applied on every classifier we have. First we select a classifier, classifier X. To calculate expected value of a combination we need to choose another classifier such as classifier Y from our pool of classifiers. Since we have 10 different Boolean operations we could produce 10 different new classifiers by combining classifier X and Y. Here, we explain our formula for one of the Boolean operation but we could easily extend it for all the rest.

As we explained earlier we could identify each classifier by its TPR-FPR on ROC curve, therefore we want to calculate expected value of TPR-FPR of the new classifier. To calculate TPR and FPR of the combination (tpr_C, fpr_C) from Classifier X with (TPR_X, FPR_X) and classifier Y with (TPR_Y, FPR_Y) we use the equations below:

$$TPR_{C} = B1 \times TPR_{X} \times TPR_{Y} + B2 \times TPR_{X} \times (1 - TPR_{Y}) + B3 \times (1 - TPR_{X}) \times TPR_{Y} + B4 \times (1 - TPR_{X}) \times (1 - TPR_{Y})$$
(3.7)

$$FPR_{C} = B1 \times FPR_{X} \times FPR_{Y} + B2 \times FPR_{X} \times (1 - FPR_{Y}) + B3 \times (1 - FPR_{X}) \times FPR_{Y} + B4 \times (1 - FPR_{X}) \times (1 - FPR_{Y}))$$
(3.8)

Which B1, B2, B3, B4 is binary representation of our Boolean operation (And = '1000', OR = '1110', XOR = '0110', etc.) For example for XOR we would have:

$$TPR_{C} = 0 \times TPR_{X} \times TPR_{Y} + 1 \times TPR_{X} \times (1 - TPR_{Y}) + 1 \times (1 - TPR_{X}) \times TPR_{Y} + 0 \times (1 - TPR_{X}) \times (1 - TPR_{Y}) \quad (3.9)$$

$$FPR_{C} = 0 \times FPR_{X} \times FPR_{Y} + 1 \times FPR_{X} \times (1 - FPR_{Y}) + 1 \times (1 - FPR_{X}) \times FPR_{Y} + 0 \times (1 - FPR_{X}) \times (1 - FPR_{Y})) \quad (3.10)$$

Now we have 3 classifiers X, Y and the combination of them using XOR operation (Classifier C) as it's shown in the picture below. Now we need to evaluate the goodness of classifier C compared to X and Y. Area Under the Curve (AUC) is a function, which can be used to evaluate the goodness. Based on the coordination of each classifier on ROC curve we can compute AUC of the classifier with the following formula:



Figure 3.7: Combination of classifier X and Y lead to C

We compute AUC for classifier X, Y and C and keep classifier C only if its AUC is greater than AUC of both classifiers X and Y. Greater AUC means that our classifier is much closer to the ground truth (which will give us the TPR = 1 and FPR

= 0). If AUC_C is greater than AUC_X and AUC_Y we keep the difference in AUC as the improvement of that classifier with the specific Boolean operation used. The variable that keeps the improvement for classifier X is IMP as defined below: $IMP_X =$ $AUC_C - Max(AUC_X, AUC_Y)$ We calculate IMP_X for all 10 Boolean operations and add them to the previous value. At the end, IMP_X has the expected value of all of the combination possible between classifiers X and Y using 10 Boolean operations. We repeat these steps for classifier X in combination with the rest of classifiers, choosing one each time as a new classifier Y and update IMP_X . The final IMP_X value has the total expected value of the improvement with respect to all the combinations if we select classifier X. After we compute IMP_X for classifier X and the rest of classifiers we will update the value in our table in a $TPR_X FPR_X$ position with IMP_X . We repeat all these steps for every classifier we have and update their IMP in the table and then we can create a table shown below with each cell representing the expected value of the classifier combinations getting improved. We select that classifier with the most improvement.



Figure 3.8: expected value of all classifiers in the ROC space

As shown in Figure 3.8 the representation of that table if we use color plotting. The red region shows classifiers with higher IMP and the blue region shows classifier with low IMP. That means if we select our basket of classifiers from red region we would have a higher chance to get better classifiers after combination of those pruned classifiers. Our assumption for creating above table was that we have all the possible classifiers with 2 digits precision and after computing all possible combination for all the classifiers we set the value of each cell below minor diameter to zero (all the classifiers below random line) since want to evaluate goodness of classifiers better than random guess. We can re-produce the same table by computing the expected value of combination just for the classifiers that already we have; by doing that we will get the below table (plot) and we can prune our classifier based on this new table.

3.3 Complexity Analysis

Given K soft detectors, let n_i be the number of decision thresholds or crisp detectors produced by each of the soft detector D_i , i = 1, ..., K, on the validation set \mathcal{V} . Let $n = \sum_{i=1}^{K} n_i$ the total number of crisp detectors in the ensembles, and $n_{avg} = n/K$ the average number of crisp detectors produced by soft detectors.

A brute-force search for optimal combination is infeasible in practice due to the doubly exponential combinations. In fact, for n crisp detectors there are 2^n possible outcomes that can be combined in 2^{2^n} ways, which makes the brute-force combination impractical even for small n values [24]. Even only pairwise combination of n crisp detectors, which requires $O(n^2)$ Boolean operations, may not be feasible in practise for large n values. The sequential combination of the IBC algorithm reduces its worst-case time complexity to $O(n_{avg}^2 + Kn_{avg})$ Boolean operations.

Both pruning techniques are capable of reducing the size of the selected subset of detectors (for Boolean combination) up to a user defined maximum number (U). The worst-case time complexity required by MinMax-Kappa technique to select Ucrisp detectors (and prune the rest) is $O(n(\log n+1)+U^2)$. It requires about $n(\log n+1)$ operations for computing and sorting the Kappa values for all crisp detectors, and U^2 for the pairwise Boolean combinations of the U retained detectors. For ROCCH-Kappa technique however the worst-case time complexity is of the order $O(n(\log n + 1))$ n_{ev}) + U^2), where n_{ev} is the number of emerging vertices which is typically around ten. The additional n_{ev} factor is due to the computation of Kappa is repeated n_{ev} times for each emerging point on the ROCCH. For numerical comparison, Table 4.2 shows the average time for each combination and pruning technique used in our experiments.



(b) κ -tpr diagram

Figure 3.4: Illustration of selected detectors (large blue circles) based on MinMax-Kappa pruning technique. All remaining detectors (small black dots) are pruned.



(b) κ -tpr diagram

Figure 3.5: Illustration of selected detectors (large blue circles) based on ROCCH-Kappa pruning technique. All remaining detectors (small black dots) are pruned.



Figure 3.6: Illustration of the selected detectors for combination mapped onto the ROC space (large blue circles). All other detectors (small black dots) can be pruned.

Chapter 4

Experiment and Results

4.1 Experiments on System Call Data set

We evaluate the accuracy and efficiency of our pruning techniques using a modern system call data set, called ADFA Linux Data set (ADFA-LD), which has been recently made publicly available on the website of the University of New South Wales [26]. The ADFA-LD data set is generated by exploiting various security vulnerabilities in a Ubuntu operating system (OS) hosting a web server. The systems consists of a fully patched Ubuntu Linux 11.04 OS with an Apache 2.2.17 web server, PHP 5.3.5 server side scripting engine, TikiWiki 8.1 content management system, FTP server, MySQL 14.14 database management system and an SSH server. First they collect normal system call traces by letting users perform basic operations, such as web browsing and Latex document preparations under controlled situations. The anomalous system call traces are collected while the system is being under six types of attack vectors resulting in a total of 60 attacks. These attacks were launched by a certified penetration tester against the system and included web-based exploitation, simulated social engineering, poisoned executable, remotely triggered vulnerabilities, remote password brute-force attacks and system manipulation. The authors of the ADFA-LD have organized the data set into 833 normal traces for training the anomaly detectors, and 4373 normal traces and 60 anomalous traces for testing.

In our experiments, we used the following experimental setup. First, we trained 20 HMMs with different numbers of states (i.e., K = 20 soft detectors), using the 833 normal traces provided for training in the ADFA-LD data set. On average, the output of each HMM is thresholded into $n_{avg} = 100$ thresholds or crisp detectors, which provides $n = K \times n_{avg} = 20 \times 100 = 2000$ crisp detectors in the ensemble.

The objective is to choose the most accurate subset from this ensemble for combination while pruning the remaining detectors according to our MinMax-Kappa or ROCCH-Kappa technique. Then, the ROC curves and AUC results are compared to those of IBC and the pairwise brute-force Boolean Crisp combination (BBC2). BBC2 is another baseline, which is the PBC Algorithm 2, but without any pruning mechanism.

For evaluation of performance, a 5-fold cross-validation (5FCV) is applied to the test set comprising the 4373 normal and the 60 anomalous traces. Since the number of anomalous traces is relatively small with reference to the normal ones, we applied the cross validation to partition the normal and anomalous sets separately, such that we keep the same ratio (normal to anomalous) and guarantee that all folds include the anomalies. Accordingly, each fold contains 874 traces selected at random from the 4373 normal traces and 12 attacks traces selected at random from the 60 attack traces.

In contrast with the standard way of applying the 5FCV, we used one fold for computing the Boolean combination according to each of the combination technique, and the remaining four folds (i.e., 3498 normal traces and 48 attack traces) are used for evaluating and benchmarking the detection performance. This is because we wanted to test the boost in performance while using a small number of anomalies during combination. The results are shown in section 4.2. For the second data set, we used the Canali data set [49]. The data was generated from 10 different machines. Our models, which are 4 HMMs and 5 SVMs, are trained on the normal data from: anubis-good + machine1 + ... + machine9, and the scores that we used are computed on the test data: machine10 that consist of (23 normal traces) + malware (5855 anomalous traces).

With this data set, in contrast with ADFA that we showed the boost in result with a small number of anomalies we want to show even when anomalies are the majority of the test set with PBC technique there is a boost in result. Here again we applied the PBC technique with help of different distance correlation metric (Kappa, Distance Covariance, Distance Correlation) and compare the result with BBC2 and IBC. For Canali we used 5FCV approach as well. The result are shown in section 4.3.

4.2 Results (ADFA)

Figure 4.1 shows the ROC curves and the AUC performance for both pruning techniques proposed in this paper (MinMax- and ROCCH-Kappa) compared to those of IBC and BBC2 (PBC with no pruning mechanism). And Figure 4.2 shows the same ROC (the same fold) for both MinMax- and ROCCH- technique with different distance correlation as discussed in section 3.2.5.

These results are for one of the 5FCV experiments (the combinations are computed on one fold and evaluated on four) as described in Section 4.1. As shown in the figure the results are comparable both in term of AUC values and the shape of the ROC curves. This is also confirmed in Table 4.1, where the AUC values of each combination technique is averaged over the 5FCV experiments.



Figure 4.1: One of the ROC curves results of 5-folds cross-validation of Boolean combination on one fold and evaluated on four folds. The numbers represent the crisp detectors selected for combination (by each technique) to achieve about the same operational point as denoted by the large black circle

Table 4.2 shows the average over the 5FCV of the pruning time (for MinMax-Kappa, ROCCH-Kappa, MinMax-dCor, MinMax-dCov, ROCCH-dCor, ROCCH-dCov), combination time, and number of Boolean operations required by each technique to achieve the final ROCCH (as shown in Figure 4.1 and Figure 4.2 for one fold).

We set the maximum number of selected detectors to U = 50 for all the abovementioned techniques (this can be further optimized, but gave good trade-off between performance and time complexity). Therefore, the input to these pruning techniques



Figure 4.2: The same ROC curves results of 5-folds cross-validation of Boolean combination on one fold and evaluated on four folds that was used in Figure 4.1 with other distance correlation metrics

is n = 2000 crisp detectors, while the output is a subset of a maximum size of 50 detectors provided for Boolean combination (in PBC Algorithm). MinMax-Kappa took about 1.6 seconds in average to select the best subset, while ROCCH-Kappa took about ten times more due the n_{ev} factor, which is described in Section 3.3, to prune the ensemble of n = 2000 detectors. Furthermore, MinMax-Kappa is able to select a smaller number of detectors than U = 50 on average and computes the Kappa values once, which explains the reduction in combination time and number of Boolean operations compared to those of ROCCH-Kappa, as shown in Table 4.2. Thereby, although both techniques provide similar AUC performance, MinMax-Kappa is slightly preferred due to its improved efficiency compared to ROCCH-Kappa.

Table 4.2 also shows that our MinMax-Kappa techniques was able to achieve the same AUC performance of BBC2 (the pairwise Boolean combination of the 2000 detectors), by selecting less than 50 detectors out of the 2000 ones. More interestingly, MinMax-Kappa achieved these results with an average time of three magnitudes lower than that of BBC2, and about 200 times fewer Boolean operations.

Compared to IBC, MinMax-Kappa requires, on average, slightly more combination time and a larger number of Boolean operations to achieve the same AUC performance. However, the number of selected detectors and Boolean functions required to realize each vertex on the ROCCH is on average five times more according to our experiments. For instance, to achieve the final operating points denoted with a large black circle on Figure 4.1, MinMax-Kappa uses two detectors and one Boolean function according to the following formula:

$$\neg D_1 \oplus D_2$$

Note that in Figure 4.1, we only shown the number of crisp detectors not to clutter the figures (i.e., 1 means D_1)

Similarly, ROCCH-Kappa uses the following detectors and Boolean function to achieve similar operating point:

$$\neg D_3 \oplus D_4$$

However to achieve similar point of operations, IBC requires eight detectors combined according to the following Boolean operations:

$$((((((D_5 \oplus D_6) \land \neg D_{11}) \land D_7) \oplus D_8) \equiv D_9) \lor D_{10}) \land D_{12})$$

The average number of combined detectors on the final ROCCH over the 5FCV is about 10 detectors when using IBC compared to two detectors when using PBC with MinMax-Kappa pruning as shown in the last column of Table 4.2. This sequence of combination rules grows linearly with the number of soft detector K, which makes



Figure 4.3: One of the ROC curves results of 5-folds cross-validation of Boolean combination on four folds and evaluated on one fold.

IBC results difficult to analyse and understand for large K values. In contrast to IBC, the combination of two detectors according to combination of PBC with MinMax-Kappa are insensitive to order in which detectors are input to the algorithm, which makes the search for the best subset of detectors easier. However, MinMax-Kappa requires an optimization of the maximum number of detectors U that trades off the complexity and the accuracy. Setting an ADS based on two HMMs into operations, requires less time and memory resources to provide the output probabilities of the input system call sequences. In addition, operating a small number of detectors becomes critical in application of anomaly detection mobile security, due to the constraint on power resources.



Figure 4.4: The same ROC curves results of 5-folds cross-validation of Boolean combination on four folds and evaluated on one fold that was used in fig. 4.3 with other distance correlation metrics.

We conducted an alternative case study to check the impact on ROC and AUC performance of all Boolean techniques, when the number of anomalies is increased during the design phase. Therefore, instead of using one fold (comprising 12 attack traces and 874 normal traces) for selecting the crisp detectors and the corresponding Boolean functions, as described in Section 4.1, we used 4 folds (48 attack traces and 3498 normal traces) and one fold for evaluation of performance.

Figure 4.3 and figure 4.4 show the ROC curves and the AUC performance for all techniques, PBC with MinMax-Kappa, ROCCH-Kappa, MinMax-dCor, MinMaxdCov, ROCCH-dCor, ROCCH-dCov, IBC and BBC2. Again, the presented results are for one of the 5FCV experiments (but the combinations are computed on four folds and evaluated on one). As shown in the figure, all techniques provide comparable results; however, with a large improvement in detection accuracy over those presented in Figure 4.1. For instance, for detecting all attacks (tpr = 100%) the false positive rate is now $fpr \approx 2\%$ compared to $fpr \approx 16\%$ in Figure 4.1. This boost in performance can be also seen in Table 4.3 in terms of average AUC values for each combination technique over the 5FCV experiments. The results of this experiments show, as expected, that when the system is provided with more normal or attack traces, the overall performance of all Boolean combination techniques improves. In such cases, there is no need to retrain the original detectors (HMMs in our case), which is a time consuming process, but the design phase of Boolean combination techniques must be repeated. This provides an advantage for IBC and PBC, since they are efficient in selecting the detectors for final operations. However, our PBC approach will always provide two detectors for each emerging point on the ROCCH, which is less costly to operate and easier to analyse in real-world setting.

4.3 Results (Canali)

Figure 4.5 shows the ROC curves and the AUC performance for both pruning techniques proposed in this paper (MinMax- and ROCCH-Kappa) compared to those of IBC and BBC2 (PBC with no pruning mechanism). They are all applied on Canali data set in order to confirm the consistency of PBC technique regardless of the data set. Figure 4.6 demonstrates the same ROC (the same fold) for both MinMax- and ROCCH- technique with different distance correlation as discussed in section 3.2.5.

These results belong to one of the 5FCV experiments (the combinations are computed on one fold and evaluated on four) as described in Section 4.1. As shown in this figure the results are comparable both in terms of AUC values and the shape



Figure 4.5: One of the ROC curves results of 5-folds cross-validation of Boolean combination on four folds and evaluated on one fold on Canali.

of the ROC curves. This is also confirmed in Table 4.4, where the AUC values of each combination technique are averaged over the 5FCV experiments.

In this experiment; we set the maximum number of selected detectors to U = 50for all of the above-mentioned techniques similar to what we did for ADFA 4.2. Therefore, the inputs to these pruning techniques are n = 900 crisp detectors (4 different HMMS and 5 SVMs each thresholded 100 times), while the output is a subset of a maximum size of 50 detectors provided for Boolean combination (in PBC Algorithm).

Compared to ADFA's result in 4.2 here we can see that ROCCH-Kappa generates even better results than IBC, however the difference is not significant yet it



Figure 4.6: The same ROC curves results of 5-folds cross-validation of Boolean combination on one fold and evaluated on four folds that was used in Figure 4.5 with other distance correlation metrics

persuade us more to use PBC; Since by using PBC we are guaranteed that our final formula is shorter and we have much fewer number of combinations that leads to less calculation time which saves CPU cycles (and get the result very much faster).

Running PBC as discussed in section 3.2 with completely new data set such as Canali (which is for Windows) and getting promising result is a proof that PBC technique is applicable on various data sets as long as having a diverse detectors trained on the them.

The key for PBC to work well is to have as many diverse detectors as possible. And therefore the complexity of the detectors is not a determining factor and we only need to keep the diversity of the pool of detectors high. With the help of PBC we prune the pool to create the basket of detectors for combination. In conclusion, PBC with reasonable distance correlation tries to select few samples from each set of detectors such that it represents the original pool.

4.4 Threats to Validity

A threat to internal validity exists in the implementation of the IBC, BBC2 and PBC algorithms as well as in conducting the experiments for anomaly detection. We have mitigated this threat by manually verifying the outputs.

We have conducted experiments using only one system call data set derived from the Linux operating system, which consists a threat to external validity of this study. More experiments are therefore required to generalize the presented results to other data sets, operating systems and other software vulnerabilities.

Evasion attacks could also pose a threat to validity. For instance, *mimicry* attacks try to mimic the normal system behavior will go undetected with the ADSs that are based on individual system calls or their temporal order [50]. Mimicry attacks could be conducted by an attacker who is able to launch his attack without tempering the normal order of system calls by, for instance, replacing foreign system call sequences (which can be easily detected) with normal ones or by using system call arguments [50].

The manifestations of such mimicry attacks could be detected by including additional features, such as system call arguments [51], return values extracted from the call stack information [52], and the user identity [53]. An added advantage of multiple detector systems, combined without our PBC, is that they can combine different detectors trained on various features, such as system call arguments, return values and other information flow features to help mitigating such evasion attacks.

Method Name	Mean	Std
BBC2	0.97426	0.001
IBC	0.97276	0.001
MinMax-Kappa	0.97074	0.003
ROCCH-Kappa	0.97051	0.003
EVBC	0.94472	0
MinMax-dCor	0.9673	0.006
ROCCH-dCor	0.97067	0.005
MinMax-dCov	0.96515	0.004

Table 4.1: Average AUC values and their standard deviations over the 5FCV for each techniques. Design on one fold and evaluated on four folds.
Table 4.2: Comparison of pruning and combination time (seconds) and number of Boolean operations required to achieve the final ROCCH during the design phase, and the number of selected detectors during the operational phase. All values are averaged over 5FCV.

		Design Phase		Operations
Method	Pruning	Combination	# Boolean	# Combined
Name	Time	Time	operations	detectors
BBC2	N/A	16364	4,000,000	2
IBC	N/A	11	11,000	11
MinMax-Kappa	1.6	15	19,701	2
ROCCH-Kappa	11.8	38	37,701	2
EVBC	370	13		2
MinMax-dCor	74	16	19,701	2
ROCCH-dCor	190	30	37,701	2
MinMax-dCov	64	17	19,701	2
ROCCH-dCov	185	37	37,701	2

Method Name	Mean	Std
BBC2	0.98177	0.00636
IBC	0.98003	0.01127
MinMax-Kappa	0.97806	0.00640
ROCCH-Kappa	0.97578	0.00585
EVBC	0.94472	0
MinMax-dCor	0.9725	0.012
ROCCH-dCor	0.9752	0.005
MinMax-dCov	0.96268	0.011
ROCCH-dCov	0.9753	0.005

Table 4.3: Average AUC values and their standard deviations over the 5FCV for each techniques. Design on four folds and evaluated on one fold.

Method Name	Mean	Std
BBC2	0.97398	0.008
IBC	0.97141	0.006
MinMax-Kappa	0.95884	0.005
ROCCH-Kappa	0.9716	0.008
EVBC	0.95074	0.006
MinMax-dCor	0.95426	0.006
ROCCH-dCor	0.96451	0.006
MinMax-dCov	0.9616	0.011
ROCCH-dCov	0.96451	0.006

Table 4.4: Average AUC values and their standard deviations over the 5FCV for each techniques. Design on four folds and evaluated on one fold on Canali.

Chapter 5

Conclusion

In this thesis, we proposed PBC, an efficient approach for selecting and combining anomaly detectors, which relies on novel pruning techniques. Our pruning technique is based on measure of agreement between detectors (here, we used Kappa, Distance Covariance and Distance Correlation).

The main advantage in the design phase is that, the PBC is able to select a small subset of diverse and accurate detectors for Boolean combinations, while discarding the remaining ones. The PBC's goal is to make a basket of detectors that complements each other's errors while still being accurate and meaningful.

Our pruning techniques we developed at the core of PBC rely on maximum and minimum of each metrics (MinMax-Kappa, MinMac-dCor, MinMax-dCov) and on the ROC convex hull (ROCCH-Kappa, ROCCH-dCor, ROCCH-dCov). This approach enables our technique to aggressively prune redundant and trivial detectors.

The results on ADFA-LD system call data sets and Canali show that PBC with both pruning techniques are capable of maintaining similar overall accuracy as measured by the ROC curves to that of IBC and BBC2 which we used for baseline cases. Therefore, our proposed PBC-based ADS is able to prune and combine large number of detectors without suffering from the exponential explosion in number of combinations provided with the pairwise brute-force Boolean combination techniques. This has been shown analytically (in the time complexity analysis) and confirmed in the experimental results.

The core of our pruning technique is choosing a good measure of agreement. Our experiments show that if we use a good and strong metric in combination with either MinMax or ROCCH techniques; the PBC works well and leads to promising results.

During the operational phase, PBC with both pruning techniques only provides two crisp detectors for each combination, while IBC requires an average of 11 detectors to achieve the same operating point (in terms of true and false positive rates).

In addition to above advantages, the proposed PBC approach is also versatile and can be applied to combine any soft or crisp detectors or two-class classifiers in a wide range of applications that requires combination of decisions.

5.1 Future Research Opportunities

For future work, we can extend our pruning technique to many different levels. Here we mainly focused only on pruning and combining in the decision level; however, there is plenty of opportunities to investigate and incorporate this technique in different levels; especially on the matching score level as you can see in Figure 3.1.

Another interesting direction is to investigate the potential improvements made by re-combining the resulting combinations with the selected detectors. We can also perform more iteration and find a solution to optimize the maximum number of iterations that we apply on our data set. We recommend to explore other measures of diversity to determine whether there is one good solution for every data set. We also intend to implement this new techniques in TotalADS [54], a tool we have developed to support multiple anomaly detectors. Finally, we need to investigate how we can reduce the size of traces to enable better scalability. Example of trace abstraction techniques are presented in [55, 56].

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